

# PKI Scalability Issues

Adam J Slagell  
slagell@ncsa.uiuc.edu

Rafael Bonilla  
bonilla@ncsa.uiuc.edu

## 1 Introduction

We spend much of our time trying to communicate with each other. Wide-spread use of the Internet has increased the number of ways and the amount we communicate with each other. For example, we may now spend many hours per day simply writing and replying to e-mails. Like normal communication, there is information that can be publicly known (or at least we do not care if someone else knows it), and there are critical messages that we prefer only to be in the possession of the intended recipient(s). Going into a dark alley to send an e-mail does not mean that it was delivered to the right person and that the information remains confidential.

Encryption helps solve problems of confidentiality. Private key or symmetric encryption systems transform, by applying complex mathematical functions, our secret message written in plain language to something that will look like gibberish. In order to reverse the transformation you need to know the correct key. Any two users trying to communicate securely can agree on a shared secret key and use symmetric encryption systems to protect their information. If the same user wants to talk with a third user, they need to agree on another key. In the end, symmetric encryption systems that wish to support communication between members of arbitrary subgroups need  $\Theta(n^2)$  private keys, where  $n$  is the number of users (one key for each pair of users). This is quite impractical.

Public key or asymmetric cryptography allows individuals to define two keys: a public one for encryption, and a private one for decryption. Now, instead of agreeing on one private key, Alice can encrypt a message for Bob using his public key and send it. Bob, knowing the corresponding private key will decrypt the message and read it. Eve, a malicious user listening to Alice and Bob's communications, will not be able decrypt the messages because she does not know the private keys. The total number of secret keys per user is reduced from  $n - 1$  to just 1. A problem with asymmetric encryption is that it is significantly slower than symmetric encryption systems. We can solve this problem using asymmetric systems to agree on a per-session symmetric key to be used for the bulk of the encryption workload. Still, a major problem remains. How can be Alice sure that the key is actually Bob's public key and not Eve's public key?

*Public Key Infrastructures* (PKIs) help solve this problem. The purpose of a PKI is two-fold: (1) to help Alice retrieve Bob's public key and (2) to give Alice confidence that the key really belongs to Bob. There are several PKI implementations. The lack of standards and the need to have a solution that can be easily used, even for large environments, have delayed the global adoption of a PKI. We have been studying the scalability of PKIs, and in this report we present several current PKI implementations and discuss the most important issues related to them. <sup>1</sup>

In section 2 we present an overview discussion of different PKIs. Section 3 describes different problems with traditional PKIs during enrollment and certificate issuance along with three different

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PKI solutions to those problems. Section 4 discusses several certificate revocation systems and discusses scalability issues with each. People are now trying to enhance PKI by providing real-time services; section 5 reviews some of those services. Then, in section 6 we discuss PKI issues that are of special interest to military scenarios. Section 7 presents conclusions and future work that can be added to our study.

## 2 PKI Overview

First, we review *Public Key Infrastructure using X.509* (PKIX), one of the two most popular PKIs. PKIX is based on the ITU-T Recommendation *X.509 Public Key Certificates* (PKC), and its study will help us better understand directory-based PKI solutions. Then we provide an overview of *Simple Public Key Infrastructure* (SPKI) which is an effort to produce a certificate structure and operating procedure that is easy to use, simple and extensible. We conclude this section with a brief discussion of *Pretty Good Privacy* (PGP), the other of the two most popular PKIs. PGP bases its structure on a so called web-of-trust where users decide which keys must be trusted and at what levels.

### 2.1 PKIX

PKIX is a PKI that uses X.509. The X.509 standard specifies a certificate format and procedures for distributing public keys via PKCs signed by *Certificate Authorities* (CAs). PKIX defines the PKI system architecture along with an X.509 PKC profile and standard procedures for registration, initialization, certification, key generation, recovery, update, expiration and compromise, cross-certification and revocation of certificates.

The architectural model consists of five components as specified in [1]:

- CAs that issue and revoke PKCs.
- *Registration Authorities* (RAs) that vouch for the binding between public keys and certificate holder identities or other attributes.
- PKC owners that can sign digital documents and decrypt documents using private keys.
- Clients that validate digital signatures and their certification paths from a known public key of a trusted CA and encrypt documents using public keys from certificates of PKC holders.
- Repositories that store and make available PKCs and *Certificate Revocation Lists* (CRLs).

In order for an individual to start using the PKI, she first needs to register by sending a request for a PKC to a CA. Along with the request, users must provide some other information like name (e.g., common name, domain name, IP address), and some other attributes to be put in the PKC. Prior to the creation of a certificate, the CA must verify that the information provided by the user is correct and that the name belongs to that user. This process of verification can be done directly by the CA, but it is more commonly done by RAs. An RA can verify the identity of the user at the moment it receives a request for a PKC, and then it will forward the request and the verified information to the CA which will create the certificate, sign it with the CA's private key, and distribute it to the user. The main idea of a certificate is to bind an identity with a public key. The public-private key pair can be generated by the CA or the public key can be presented by the user as part of the attributes. If the key pair is generated by the CA, then it must be sent back

to the user by trusted means. If the user provides the public key, she must prove that she has the corresponding private key.

The initialization process consists of an *End Entity* (EE) (e.g., a client using a web browser) retrieving all the values needed to start communicating with the PKI, like the CA's public key that will enable the subject to verify PKCs signed by the CA. If Alice wants to communicate with Bob (not with someone else claiming to be Bob), she must first go to a repository and retrieve Bob's certificate. These repositories are like phone directories with certificates indexed by users names. One difficulty at this point, called the *John Wilson Problem*, is how Alice can be sure that she has the correct John's certificate and not some other John's certificate. A partial solution is achieved by having the CAs verify the names during enrollment to assure they are locally unique. Additional information could be added to the certificate's name so it will be different from all other names issued by that one CA. But we can still find two (and probably more) John Wilson with certificates issued by two separate CAs.

Once Alice has Bob's certificate signed by some CA, she can verify it if she trusts the CA and has already its public key. If not she has two options: discard Bob's certificate or get Bob's CA's certificate. After verifying the certificate(s), she can use Bob's public key. Now she can communicate securely with Bob by encrypting messages using his public key. These messages can be part of a session key sharing protocol, such as in [10], in order to use faster symmetric key cryptography for the remaining communications. For further proof of identity, Alice can send a challenge to Bob encrypted with his public key. Only Bob, knowing the corresponding private key, will be able to decrypt the challenge and respond to it, thus proving his identity.

Key pairs need to be updated regularly and new PKCs issued mainly for two reasons: the key pair has exceeded its predefined lifetime or the private key has been lost or compromised. In either case, the PKI must provide a smooth transition from the old key pair to the new one. The worst scenario is when the root CA's key has been compromised. In this case, the root CA must generate a new key pair making useless the paths underneath it in the hierarchy until all the revoked certificates issued by the root CA are replaced with new PKCs. X.509 defines one method to revoke certificates where each CA periodically (e.g., hourly, daily, or weekly) issues a signed list containing the serial numbers of revoked certificates called a *Certificate Revocation List* (CRL). Besides checking the signature of the certificate, clients should get a recent CRL and check that the certificate is not in the list.

One more aspect defined in PKIX is *cross-certification*. Cross-certification is used to allow users under one CA or domain to communicate securely with users under a different CA or domain when the CAs do not share a common root. Cross-certificates can be issued in one direction or in both directions between two CA's.

The PKCs we have discussed so far are used to perform identity-based access, but for many systems rule-based or role-based access is desired instead. These forms of control require additional information that is not normally included in PKCs. PKIX defines an *Attribute Certificate* (AC) that binds this extra information as a digitally signed data structure with a reference back to a specific identity based PKC or to multiple such PKCs. Separating identity certificates from attribute certificates is good practice because attributes/roles change frequently while identities remain constant. *Privilege Management Infrastructure* (PMI) is defined in [1] as the set of hardware, software, people, policies and procedures needed to create, manage, store, distribute, and revoke ACs.

## 2.2 SPKI

Simple Public Key Infrastructure [4] is aimed to provide an easy, simple and extensible form of PKI with the main purpose being authorization rather than identification. SPKI defines authorization certificates in addition to identity certificates used by PKIX. Certificates come in three categories: *identity certificates* which bind a name to a key <name,key>, *attribute certificates* which bind an authorization to a name <authorization, name>, and *authorization certificates* which bind an authorization directly to a key <authorization, key>.

The John Wilson problem in section 2.1 proves that names cannot always work as identifiers and this is a serious drawback of PKIX. CAs already have to locally distinguish between John Wilsons. To globally extend names, users need to know the extra information added to locally distinguish and the issuing CA, in order to identify the correct John-Wilson's certificate. In contrast, SPKI uses *Simple Distributed Security Infrastructure* (SDSI) names to create globally unique identifiers. An SDSI name is an S-expression with the word "name" and the intended name. For example, jim: (name rafael) is the basic name "rafael" in the space defined by jim. SDSI names can also be compound, for example, jim: (name rafael adam) is the basic name "adam" defined by rafael and indirectly referenced by jim. There are several ways to make names globally unique identifiers. Because keys -and most likely their hashes- are unique they can serve as unique identifiers. Fully-qualified SDSI names must include the name of the space in which they are defined. SPKI supports compatibility with X.509 names by converting those names to SDSI names, for example (name <root key> <leaf name>) and (name <root key> <CA1> <CA2> ... <CAk> <leaf name>) are examples of X.509 names converted to SDSI names.

The authorization process can be summarized in 6 steps:

1. Alice wants to access a resource and asks the resource owner (or administrator) to grant her access.
2. The owner decides if the request is valid and what level of access should be granted to Alice.
3. The owner creates an authorization certificate for Alice binding a public key, for which Alice has the corresponding private key, to an ACL and signs it. The certificate must be sent back to Alice.
4. Alice presents a signed request to access the resource. Alice's authorization certificate accompanies this request.
5. The resource manager checks that the authorization certificate is valid (i.e., signed by the resource owner) and confirms that the signature was made by the key in the certificate.
6. Finally, if either the certificates is invalid or the signature is bad, the request is denied. Otherwise, Alice gains access to the resource.

As an alternative, authorization could have been performed using a combination of identity and attribute certificates as in PMI. Here, Alice can have an SDSI name bound to a public key by an identity certificate, and an attribute certificate binding an authorization to her identity. The identity in both certificates acts as a mapping field. Alice must present both certificates when asking for access. The resource manager can check the authorization in the attribute certificate as before but also checks the identity certificate looking for a match with the identity specified in the AC. If the authorization is correct and the identities match, access is granted. This has the benefit of being more easily audited. However, anonymity may be preferred in some cases. This is a goal that cannot be met using attribute certificates.

Two more aspects of SPKI are delegation and threshold certificates. Authorization certificates can give users the power to delegate authorization to another user without having to ask for a new certificate from the owner of the resource. Delegation can be in full or limited by the delegator. Threshold certificates are defined by splitting the right of access between  $n$  subjects and specifying a threshold value  $k$ . The authorization process now works by having  $k$  subjects present a request for access. Only when the threshold value is met can access be granted.

Validation and revocation of certificates under SPKI, as in PKIX, is handled by time-constraining certificates with not-before-dates and not-after-dates and by using CRLs. Upon receiving an SPKI certificate, the validity period is checked, and then the certificate's serial number is compared against those in the most recent CRL.

## 2.3 PGP

Pretty Good Privacy (PGP) was designed by Phil Zimmermann in 1991. PGP differs completely from PKIX in its distributed approach to key management. PGP does not use certificates and registration authorities. Instead, PGP implements the concept of a "web-of-trust" where users generate their key pairs, distribute their public keys and ask other PGP users to sign their public keys, thus constructing a web of users trusting each other.

Alice, a business representative attending a conference in Boston, meets Bob, a business consultant, and after talking they realize that there are some projects in which both are interested. They decide to keep in contact, and at the end of the conference they exchange keys to securely communicate with each other. Their keys (or hashes) may be impressed in their business cards and available at some web site or directory from which they can be fetched. Carol, an acquaintance of Alice, decides to take part in these projects but wants to communicate with Bob first. Bob sends his public key to Carol but she has no way to be sure that the key is really Bob's key and not that of an impostor trying to steal from Carol, except that Bob sends his key signed by Alice (and possibly some other users). Since Carol knows Alice and trusts her to sign keys, she can be confident that the key is actually Bob's. From now on, Carol and Bob can communicate securely.

The main advantage of PGP is that users can manage their own keys. PGP does not need a central authority saying which keys are OK to trust and which keys have been compromised. PGP provides each user with a public-ring. A public-ring is a key repository where users can store keys they receive and assign level of trust to them. It is not clear yet how good it is to leave the decision about trustworthiness to end users instead of having a central authority that takes care of validation and verification as in PKIX. In the example above, when Carol receives Bob's key, she trusts it because it came signed by Alice. Alice's key is within Carol's public-ring and has a level of trust high enough (assigned by Carol since she personally knows Alice) to sign keys. Carol can have more keys in her public-ring that are trusted just for communication but not for signing other keys. Additionally, Carol can define her own policy so she will accept a new key only if it is signed, for example, by at least three other keys she trusts for signing. Carol can modify her public-ring and levels of trust at any moment. If a key has been compromised, she can delete it so she will not accept a message signed by that key. She can also accept, by her own risk, keys that are not signed or signed by people she does not know or trust.

Revocation is not formally addressed in PGP. If Alice's key has been compromised, she must communicate so immediately. Alice can create a revocation message saying that her key has been stolen and that nobody should trust a message signed by that key anymore. Finally, she must create a new pair of keys and distribute her new public key. The problem here is that Alice cannot be completely sure that every single user having her old key has received her revocation message. Instead, Alice could add an option field to her certificate pointing out her web page or a directory

where other users can check her key status. This solution does not scale well. PGP users have too many different places to check for keys status and they cannot be sure that the information is up-to-date.

Having given an overview of PKI, we now consider the main aspects of enrollment and certificate issuance.

### 3 Enrollment and Certificate Issuance

Enrollment and certificate issuance are two things users need to take care of before using PKIs. These processes can be as long and complex as in PKIX or very easy as in PGP. In this part we will refer to many of the concepts already described in sections 2.1 and 2.3 to compare both methods, highlighting some of their individual problems.

In trusted third-party methods of key management, like PKIX (Public Key Infrastructure using X.509 standard), when a user Bob wants to obtain a certificate to prove his identity, he must send a request for a certificate to the CA (a central trusted third-party). The request may contain Bob's public key or the CA may instead generate a key pair for Bob and distribute it along with his certificate. To process the request the CA must verify Bob's identity and that the public key belongs to him. After that, it will create a certificate for Bob and sign it using the CA's private key. Finally, the CA sends the certificate (and possibly the new private key) to Bob. This process sounds simple but has several difficulties as stated in [18]:

- It is hard to determine the level of trust in Bob's identity implied by his certificate.
- It is hard to define the relationship between Bob and the CA that certified his public key and to specify the relationship in his certificate.
- Having a "single trusted entity" creates security, administrative and possible legal problems.
- Certificates and keys must be securely distributed to end users and subordinate CAs.

Referral methods such as PGP solve many of the problems mentioned above. PGP employs the concept of "introducers". Introducers are users of the system signing keys of other users, presumably friends or people they know and with whom they exchange keys face-to-face. If Alice knows Bob, she can sign his key, and then when Bob tries to communicate with Carol he will present his key signed by Alice. If Carol also knows Alice, she will trust Bob's identity. This process allows users to construct a web-of-trust. Additionally, users can assign levels of trust to the keys they use; some keys may be trusted to sign other keys, and some keys may be trusted just to identify their owners. But referral methods are not a complete solution and suffer from problems like the following:

- An introducer must be sure of Bob's identity and that the public key presented belongs to him. In our example this likely means Alice has met Bob in person to get his key or at least its fingerprint.
- It is possible that Carol does not know Alice, and so she will not trust Bob.
- Currently, key revocation is not formally addressed for referral methods.

As seen above, enrollment and certificate issuance in PKIX is a process that can take a long time to finish. Online CAs enhance this process by making it faster. Online enrollment follows almost the same steps as before, but now instead of having the CA carefully verifying users information,

an online CA challenges a user with an e-mail sent to the address provided within the request. Once the user successfully answers it, the online CA will send her certificate (and maybe her private key). This method allows e-mail addresses to be bound to public keys, though it relies on the non-existent security of e-mail protocols. A more secure example of an online CA is a Kerberos CA. Here Kerberos identities are bound to keys and the identity is securely verified with a Kerberos ticket. Another difference with traditional enrollment is that online CAs usually issue short-lived certificates. Near the expiration of certificates, users may ask for new ones if needed. Unlike traditional and online enrollment, PGP provides a completely different solution. There are not central authorities that take care of the process and the certificates. Users create their own keys and start using them. Certificates gain value by the signatures of introducers. Additionally, users can publish their public key in directories where other users can retrieve them in order to communicate with each other, but this is not a requirement.

Scaling enrollment and certificate issuance presents new challenges. For PKIX, cross-certification and *Bridge CAs* (BCA), as described in [19], can be used to allow users under different domains (and possibly different CAs) to communicate with each other. The problem is that solving organizational issues (especially about the meaning of “trust”) is not always easy. We discuss this further in section 5.2. Besides that, implementing a large scale PKIX system incurs several costs. Certificate requests must be manually verified and processed; so new staff must be hired for this task. Online CAs may reduce these costs but more computational processing and good channels of communication are required. Assigning the verification process to already existing staff can be another option. Hardware related costs are also important. Those costs may be by far the most expensive if the PKI is implemented using some kind of device like smart cards to protect user’s keys. The PKI then needs to provide its users with special hardware like smart card readers as well. It can be argued that this is not a direct PKI cost but a cost for users of the PKI solution. Hardware costs must also include central equipment to work with smart cards when issuing keys and certificates. Certificate revocation and CRL distribution costs must also be considered. As indicated by the *National Institute of Standards and Technology* (NIST) in [16], a PKI should expect to revoke about 5 percent of all certificates issued each year because the corresponding private keys have been lost or compromised. Another 5 percent of certificates are expected to be revoked because of users leaving the system. One must also account for certificates generated for completely new users. It is expected that 5 percent of the certificates held in a given year will be for these new users. In contrast, the distributed nature of PGP and its no enrollment solution helps with some scalability issues, but now revocation becomes more difficult.

PKI literature presents several other works that try to improve enrollment and distribution of certificates and keys. We describe three such works and the problems each one solves.

### 3.1 FreeICP

*FreeICP* [7] combines directory methods with referral methods by having a CA hierarchy that mimics PGP’s web-of-trust model using a collaborative web-based trust scoring system. FreeICP proposes a CA hierarchy with a root CA that certifies two types of intermediate CAs: *Entry Level* (EL) CAs and *Verified Identity* (VI) CAs. The main role of an EL CA is to generate short-lived certificates online to any user requesting one. The EL CA performs minimal validation by following a naming policy, avoiding duplicated entries and verifying the validity of the e-mail address by sending a message to it. Through EL CAs, FreeICP puts a valid, working certificate into the user’s applications immediately and for free. VI CAs issue long-lived certificates once users have met specific levels (scoring) of credibility and trustworthiness. The hierarchy can even define several CAs, each with successively more stringent scoring requirements. The VI CAs also have both X.509

certificates and PGP key-pairs so they can act as cross-certifiers.

An EL CA certificate gives the user a fully-functional way to identify herself. Applications needing higher levels of trustworthiness can insist on a VI CA certificate, forcing the user to get one by improving her score. The scoring system consists of a policy specifying different types of proof of identity that a user can present and the points (score) assigned to them. It also specifies two types of validators that are in charge of collecting these proofs: automatic validators and user-driven introductions. Automatic validators are programs that verify some of the user's personal data through automated queries on public websites. Addresses and phone numbers, country-specific identifiers in public national databases, PGP key-based introduction, photographs and other human-verifiable data are examples of personal data collected by an automatic validator. User-driven introduction deals with FreeICP users introducing new users to the system and users presenting cross-certification from other CAs as a proof of identity. One last advantage is that the scoring process is a natural solution for contention. If two or more users are claiming the ownership of certain identity, the dispute will be solved by giving the identity to the user with the highest score since scores are improved by presenting more and better proofs of identity.

FreeICP solves the problems of:

- the level of trust to assign a user's identity by employing a scoring system to reflect trustworthiness.
- the relationship between the user and the CA and the way it is implied in the user's certificate. The CA plays an active role in the verification of the user's identity. Recall that a VI CA certificate is issued once the user has proved, with certain level of trust, his or her identity.
- not being able to control the trustworthiness of their certificates, as viewed by others, which is a problem in PGP. The scoring system allows users to improve the trustworthiness of their certificates.
- contention.

### 3.2 Self-Assembling PKI

In [5], Jon Callas presents a *Self-Assembling PKI* as a new way of constructing certificates that helps PKIs provide a widespread deployment of secure communications. Self-Assembling PKI uses existing PKIs, security standards, and systems to achieve its goals. The infrastructure consists of a server sitting within the network that creates keys and certificates for all of the network users. By sitting inside the network, the program notices the presence of already authenticated users (users using the network must have been authenticated before by another system, probably by providing a combination of user name-password) and automatically creates certificates for them. These certificates can be augmented as more information is learned about the users. Notice that no additional enrollment is necessary since the user has been already authorized to use the network, and we assume the organization owning the network has already enrolled the user and hence her identity has already been verified prior to granting access to the network.

Here is an example of the communication process described by Callas. Alice wants to securely send an e-mail to Bob. Alice connects to her usual mail server. A proxy mediates this connection, and after she successfully authenticates to the mail server, it creates a short-lived certificate for her. Alice sends the e-mail to Bob. Maybe more information is learned about Alice from this e-mail and is added to her certificate. Since Bob is a user on the same mail server, the proxy creates a short-lived certificate for him and encrypts Alice's e-mail using Bob's public key. Bob connects to his usual mail server and after successful authentication, the proxy decrypts Alice's e-mail and

presents it to Bob. As an option, the message can be modified to let Bob know that it was delivered securely.

Self-Assembling PKI provides:

- widespread deployment of secure communications.
- transparency of use.
- ease of deployment.
- risk mitigation.
- increased level of trust in users identities.
- no need for a “single trusted entity” or certification authority.
- no need for distribution of certificates and keys.
- revocation by the use of short-lived certificates.

### 3.3 The Canadian way

The work in [11] describes the concern of Canada’s Government to deliver secure online services. The main contribution of this paper is the separation of registration and enrollment for a PKI solution. Individuals will **register** with a central authority and get an *epass*. An *epass* is a pseudo-anonymous public key certificate where the identifier is a *Meaningless But Unique Number* (MBUN). At this point users are not required to identify themselves. Later on, users will need to use government programs, and they will **enroll** in such programs. The enrollment process consists of a user presenting her *epass* and proofs of identity to the program. The program will verify the user’s identity and create an association between the MBUN from the user’s *epass* with a *Program ID* (PID) number. The PID is the index for the user within the program. Enrollment must be done once for each government program on the occasion of its first use. Once enrolled, users can authenticate themselves with their *epass*, and the program will uniquely identify them by the MBUN-PID mapping.

It is interesting to notice that the Canadian way for secure online services is very similar to the ideas implemented by Microsoft in its .Net Passport single sign-on solution.

The main advantages about this idea are that:

- it provides a single sign-on solution for online services.
- data mining between organization can be done using MBUNs instead of full names as keys.
- on its own, the certificate (*epass*) contains no information about the user.
- individuals can use more than one *epass*, allowing them to fine-tune their anonymity based on their level of privacy concerns.

For our study, the Canadian way:

- increases the level of trust in user identities since each program has the ability to validate identities at its own level of concern.
- has a CA that is just an entity that issues public key certificates. Trust is now managed by each program’s identity verification process.

## 4 Certificate Revocation

Certificates are usually given a fixed lifetime, after which they expire. However, it is possible that a certificate becomes invalid before its expiration. This could happen if the private key corresponding to the certificate has been compromised. More frequently though, a person will leave a position within an organization, and the management will want to revoke the certificate to prevent them from posing as a member any further. A member could also move within an organization, thus changing the systems to which she has access. This will likely require the revocation of attribute certificates. In [14] it is estimated that 10% of certificates will actually need to be invalidated before expiration. Therefore, it is important for most PKIs to have methods to perform timely revocation of certificates. In this chapter we discuss some of those methods.

It is important to note that while most systems do have methods to deal with revocation, these can be costly to implement. Implementors of a PKI could choose not to address revocation and instead use alternatives that minimize the risk of not revoking keys. A simple solution might be to always use very short-term certificates. It takes significant time and effort to crack a key. By reducing the life of the key, the owner reduces the probability that it will be cracked while it is still valid. Another alternative is to store keys in tamper-resistant hardware. However, this only protects the private key from direct attacks. The public key is still exposed, and attacks can be mounted with just the public key information in order to reveal the private key. Of course the feasibility of such an attack depends largely upon the algorithm and key size. Additionally, tamper-resistant solutions are not based off of well understood mathematical problems that we believe to be hard; instead they are based off of electrical engineering or physics problems which have shorter lifespans. Just because something is tamper-resistant today, that does not give one confidence that it will be in a few years. For example, many tamper resistant technologies, including smart cards, have fallen prey to attacks that analyze electrical signals. This being said, we feel that tamper-resistant hardware is a good second layer of defense but should not solely be relied upon.

### 4.1 Certificate Revocation Lists

*Certificate Revocation Lists* (CRLs) were one of the first methods to revoke certificates. These so called “black-lists” are lists of all currently valid (meaning non-expired) but revoked certificates. A CA would issue one CRL for all certificates that it had revoked. In [16] it was suggested that CAs should issue CRLs on the order of every two weeks. No matter how often the CRLs are updated, it must be done in a manner that a user can verify that she has the latest CRL. This could mean that the user knows it is updated at a specific interval, or the CRLs could indicate when the next one would be issued.

Of course one of the main disadvantages of CRLs is unscalability. These lists can become quite large for a user to download. The problem is exacerbated if revocation information needs to be very fresh. In this case the CRL must be updated more frequently, and hence downloaded all the more frequently. So there is this trade-off that we often find between freshness and scalability. At the one end we could have no CRLs which is very scalable, but the information about certificates is stale. On the other end we could update daily, but this is not very scalable if a user must download many CRLs, even in the age of the networked computer.

It is important to realize that downloads need not be synchronous, though. This fact can be leveraged to provide scalability by downloading CRLs in times of low network use, such as during the evenings. Going a step further, clients could be configured to download CRLs at random times during the evening to avoid bursts of traffic. This would be better than everyone trying to download certificates at, say, midnight. It has been suggested by some to over-issue CRLs to avoid the bursts

of traffic near the release of new CRLs. Over-issuing means that new CRLs are released before all the older ones expire so that there are many different non-expired CRLs at a given moment. In [8], Cooper models how over-issuing affects the peak request rate for CRLs. While he shows that it does reduce the peak rate effectively, it is important to realize that average workload for a CA is increased and the average request rate for a directory is unchanged.

#### 4.1.1 Delta-CRLs

One of the first solutions to address the scalability problems of CRLs were *delta-CRLs*. A delta-CRL is just a list of changes to a base CRL. In this situation a complete CRL is issued regularly, but infrequently. In between issues of the base CRLs, delta-CRLs are issued that specify new revocations that have occurred since the release of the last base CRL. This reduces the amount of information that a client must download on a regular basis while still providing information that is fairly fresh. The end user must still have a mechanism to know that the delta-CRL is the freshest out there. So the delta-CRL should be issued at regular intervals, as well. The most significant disadvantage is that they still do not provide a succinct proof of validity that an end user can send to another end user with her certificate. The end user would have to store the base CRL and delta-CRL with their certificate to provide proof to an offline agent. Some new methods of revocation provide more succinct proof that a certificate has not been revoked.

In [6], Adams et al. make two improvements to traditional CRLs as discussed above. The first improvement is almost functionally identical to delta-CRLs and is more of a political difference. There is always a balance between freshness and cost in revocation systems. They feel that because not everyone may be interested in the absolutely freshest information, it makes sense to charge a premium for the freshest updates. They propose using an X.509 extension field for what they call the *Freshest Revocation Info Pointer* (FRIP). This is just a pointer to a special type of delta-CRL that contains the absolutely freshest information. This *Freshest delta-CRL* (FCRL) must be served from a trusted source now since it is issued irregularly, and the client must be assured it is the latest available. But since it is assumed that the user is purchasing the list, the server must be trusted to some extent anyway. The purchase price should be enough to make up for the cost of the CA setting up extra servers.

The user is not really benefiting from this system, except that the FCRL is more current than a regular delta-CRL. The client is still downloading as much information as she would with traditional delta-CRLs. The directory is doing less work. It is not handling the delta-CRLs at all. The CA does more work now that it must create more updates, since FCRLs contain the absolutely freshest information. Moreover, it must serve this data or rely on some trusted system to serve it. The total network traffic may decrease if people are not willing to pay for the FCRLs. It would be interesting to see if this system would work socially. People do not like to pay for something that they got for free before or things that they do not understand. If the CAs cannot get enough subscribers, the costs per user for FCRL access would be too high for most individuals. Overall, the main advantage is that the FCRL can contain the absolutely most current information, but this comes at the cost requiring a trusted server that is always online. This FCRL server can then become a point of DoS attack. Replication thus becomes necessary for resilience, but replication among several non-trusted directories is easier than replicating servers providing private data that is being sold. And what is keeping an organization from caching a very current FCRL for all its members? Now the members have fresh information at almost no cost.

Another variant of delta-CRLs is called *Sliding Window Delta-CRLs*. Presented in [9], Cooper shows how to lower the request rate of base-CRLs and the peak bandwidth at the directory by using his improved delta-CRLs. Typically a delta-CRL lists all revoked certificates since the most recently

issued base-CRL. So the window over which the information is collected for a delta-CRL varies. He suggests using a fixed window size. For example, a base-CRL may be issued daily with delta-CRLs issued every 15 minutes. The window size could be 72 hours, meaning that a delta-CRL lists all of the certificates revoked within the past 72 hours. If a user never goes say 71 hours without validating a certificate, then she will never have to download a base-CRL again! He demonstrates that this is a great improvement over traditional delta-CRLs, and he shows how to improve peak request rates further by over issuing delta-CRLs. Of course the degree of improvement depends upon optimizing the choice of the window size for the given base-CRL and delta-CRL periods.

#### 4.1.2 CRL Distribution Points

Another improvement to CRLs was specified in the X.509 v2 CRL specifications [6]. In the version 2 CRLS, *CRL Distribution Points* (also called *Segmented CRLs*) are defined. CRL distribution points fragment the CRL into smaller parts. If these fragments are organized into logical divisions, it is likely that a user will only need to download a few fragments rather than the entire CRL. The certificate specifies which distribution point corresponds to that certificate. Distribution points can be used with delta-CRLs, as well. Here the delta-CRLs are broken into fragments -most likely along the same serial number boundaries as the base CRLs- as well. CRL distribution points do help to address the problem of scalability by reducing the amount of communication between directories and end users. However, it could happen that the fragments of the CRL do not grow uniformly. Certain distribution points could grow quite large, and the partitioning of the serial number space cannot be changed later.

The second improvement by Adams et al. in [6] addresses the problem of CRL distribution points that grow non-uniformly. They create *Redirect CRLs* (RCRLs) that sit between the end user and the CRLs. The CRL distribution pointer and FRIP now point to redirect CRLs. These redirect CRLs tell users which fragment to look at for the certificate in question. This way the serial number space can be repartitioned between CRL distribution points at any time. The problem of course is that there is now more work for the CA, and the client has an extra step of indirection involved in checking any CRL or delta-CRL. We would be surprised if the benefit outweighs this extra cost. Adams et al. provide no evidence that this non-uniform growth of CRL distribution points is actually a problem nor do they indicate how much of one it is.

## 4.2 Certificate Revocation Status

An alternative to CRLs, which are large signed statements about the status of several certificates, would be signed statements about single certificates. Instead of sending CRLs every day, the CA could send separate signed statements for *every* non-expired certificate the CA has published to the directory! It would have to send both positive and negative statements about certificate status now; otherwise an untrusted directory server could simply neglect to send a negative statement, thus leading a client to believe the certificate is valid. This isn't a problem with a CRL since the client trusts the CA to indicate *all* revoked certificates on the list. An untrustworthy directory cannot simply strip out a particular certificate from a CRL without invalidating the signature on the CRL. Thus the client only has to worry about the directory not returning the most current CRL. Dating the CRL and knowing when the next one comes out allows the client to notice such misbehavior by the directory server.

Obviously, this is not a practical solution. While it does reduce the amount of information downloaded by an end user significantly, it over-burdens the CA. The CA must not only compute orders of magnitude more signatures, it has to send much more data to the directories. This extra

data is from the signatures and the fact that information about valid and revoked certificates must both be sent. However, Micali [14] does feel that this idea has merit in that it is shifting some of the burden away from the directory-to-user communication and back to the communication between the directory and the CA. With CRLs the work-load is unbalanced, and most of the traffic is between the users and the directory. Micali takes the naive solution above further by reducing the size of the signature, and hence the data transmitted, and reducing the computational work of performing signatures. By using the light-weight signatures he proposes, signature size is reduced by about one order of magnitude to 100 bits, and the computational cost of signing is reduced orders of magnitude. He calls this system *Certificate Revocation Status* (CRS).

The light-weight signatures are created as follows. Let  $F : \{0, 1\}^{100} \rightarrow \{0, 1\}^{100}$  be a fast one-way function. For every certificate that the CA issues, it creates two private values associated with that certificate called  $Y_0$  and  $X_0$ . These are each 100 bits long. Say that the CA wants to update certificate status daily and wants certificates to last for one year before expiration. Then the CA publishes  $Y = F^{365}(Y_0)$  and  $N = F(N_0)$  as part of the certificate. On day  $i$ , the CA publishes  $Y_i = F^{365-i}(Y_0)$  if the certificate is still good. If it has been revoked it publishes  $N_0$ . The user checks  $Y_i$  by verifying that  $F^i(Y_i) \equiv Y$ . If the response is instead  $N_0$ , the user checks that  $F(N_0) \equiv N$ . The security of this signature relies completely upon the fact that  $F$  cannot be inverted easily. Note that the directory cannot trick the user in any way. If the directory responds with an older  $Y_i$ , the user will detect this. If the directory responds with  $N_0$ , the certificate must be revoked since otherwise the CA would not have released the value. All the directory can do is choose not to respond, but it could do this in any revocation system.

One side effect of this system is that every day a certificate holder can get a short proof of the validity of her certificate for that day. She can bring it with her on a smart card or some other media with her certificate to prove validity to an offline agent. This is the first system we have seen that provides succinct proof of validity to the end user. However, two issues really concern us. First, the CAs now must store private information associated with every certificate. This isn't a storage issue, but a management issue. It is much easier for a CA to protect a few very important private keys from insider compromise than it is to protect tens of thousands of pieces of confidential information. The second problem is that there is limited granularity to the system, and it is fixed once the certificate is issued. It is like they are creating one time signatures, in our example one per day for a year. The computational speed of the signature algorithm is directly proportional to the lifetime of a certificate and its granularity (period of update). It is unclear exactly how much faster these one-way functions are compared to traditional public-key signature algorithms, but eventually the cost will become unbearable if the update rate is increased enough.

### 4.3 Certificate Revocation Trees

*Certificate Revocation Trees* (CRTs) referred to in [20] are the type first introduced by Paul Kocher in 1998. The basic idea of a CRT is that revocation information is provided in the leaves of a binary hash tree, and the root of this hash tree is signed by the CA. To prove that the information a directory gives to a user is true, it provides the user with the leaf node of interest and the minimum number of node values from the rest of the tree in order to recompute the root of the tree. The user then verifies the root value against the signed root that the directory provides. Any alteration to the leaves of the tree will alter the tree's root. So as long as a strong, collision-resistant hash function is used, a directory cannot deceive the user. Also, it only has to provide proof  $O(\lg[n])$  in length, where  $n$  is the number of revoked certificates. This is much more succinct than an entire CRL, and it may be possible for the end user to carry this proof along with her on a smart card or similar device to prove the current validity of her certificate.

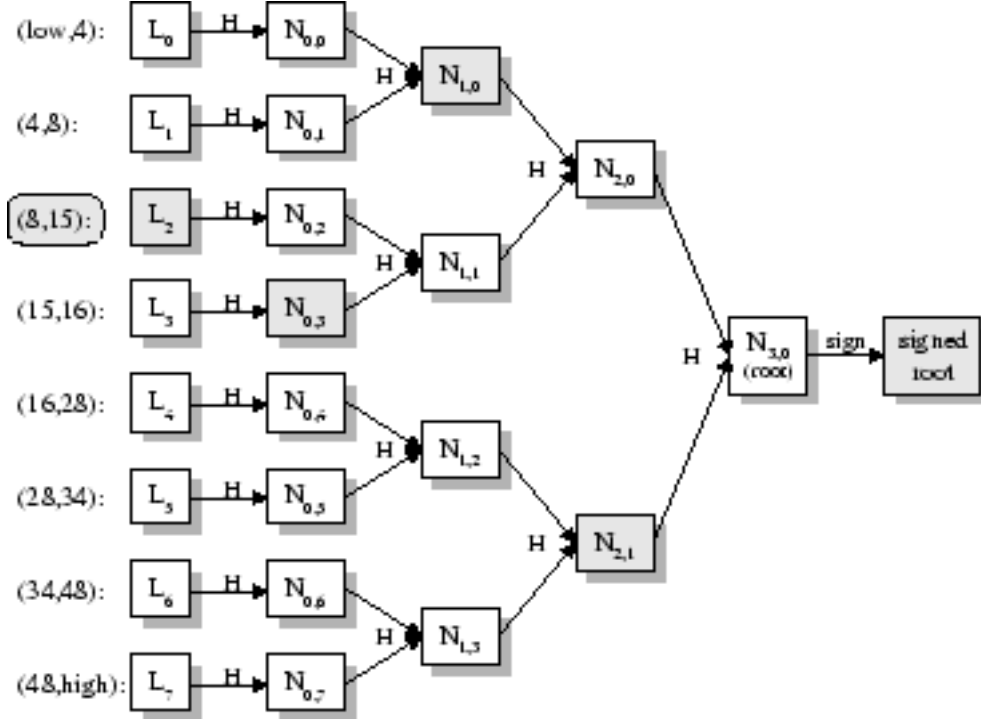


Figure 1: Example CRT from [20]

More specifically, the leaf nodes contain information of the form  $(i, j)$  where both certificate  $i$  and  $j$  are revoked, but no certificate number between them is revoked. Such a value can demonstrate that either certificate  $i$  or  $j$  is revoked, or it can be used as positive proof - for any certificate between  $i$  and  $j$  - that demonstrates validity. Consider the certificate tree in figure 1. Suppose a user queries the directory about certificate number 14. Then the directory must supply the leaf node,  $L_2$ , and also nodes  $N_{0,3}$ ,  $N_{1,0}$  and  $N_{2,1}$ . These are the siblings of all the nodes on the path from the leaf back to the root. With these nodes and the leaf, the end user can compute the root which it compares to the signed root provided by the directory.

Overall, the information sent to the directory is more than in a simple CRL. However, the benefit is that the end user needs data only on the order of a log of that which a CRL uses. This is fine since the CA is only sending data to the directory once per update, but the directory is constantly communicating data to the end users. So it makes sense to significantly reduce the data communicated with the end users, even if it comes at a small cost to the communication sent between the CA and the directory. In fact, it is only a quadratic increase in the amount of data communicated with the directory.

Naor et al. [15] improved upon Kocher's CRTs. With Kocher's CRTs it is possible that the entire hash tree must be recomputed during an update. Naor et al. sought to save this extra computational work and data transmitted to the directory by reducing the effect an update has on the hash tree. They accomplish this by using *2-3 trees* instead of simple binary hash trees. 2-3 trees have two important properties with respect to their goal: 1) membership queries, insertions and deletions only change nodes in the search path, and 2) tree nodes have bounded degree. In fact other trees with these properties could be used. They mention treaps as an alternative with their own set of advantages and disadvantages. Tree updates -removing expired certificates or adding newly revoked certificates- typically involve only the nodes on the path back to the root, but they

can also involve the addition or deletion of nodes to rebalance the tree.

In their comparisons to CRS and CRLs, Naor et al. find that they have reduced the overall communication between the CA and directory by orders of magnitude. At the same time they have kept the communication between the user and the directory small when compared to CRLs. They do not compare that communication to CRS, probably because they require more client to directory communication. They also do not compare the performance of their trees to Kocher’s. So it is difficult to predict how much of a difference their improvements make.

#### 4.4 Windowed Certificate Revocation

*Windowed Certificate Revocation* (WCR) is just an improved method of implementing CRLs, and it applies equally well to delta-CRLs. McDaniel et al. sought a balance between systems that always retrieve a fresh certificate and systems using CRLs. It is computationally costly, because of digital signatures, to always retrieve a fresh certificate, and CRLs can be costly in terms of communication, due to their large size. However, in [13] the authors should consider that always retrieving fresh certificates could be more costly than CRLs from the amortized costs of small communications. Regardless, the goal of WCR is to find a balance between the two systems through parameters chosen by the system’s users (both the certificate issuer and users of the certificate). In fact, degenerate cases of WCR turn into the above mentioned systems.

There are two main differences between CRLs and WCR. First, in WCR there must be a method for a user to retrieve a “fresh” certificate if desired. This service most likely will not be used all the time, though. WCR also maintains CRLs but with a distinct difference; certificates do not necessarily remain on the CRL until they expire. This is the second difference. A parameter called the *revocation window size* determines how long a certificate is on the revocation list. More specifically, it specifies an integral number of consecutive CRL publishing dates that the revocation information must appear on. By adjusting this parameter, the size of the CRLs can be adjusted without changing the lifetimes of the certificates.

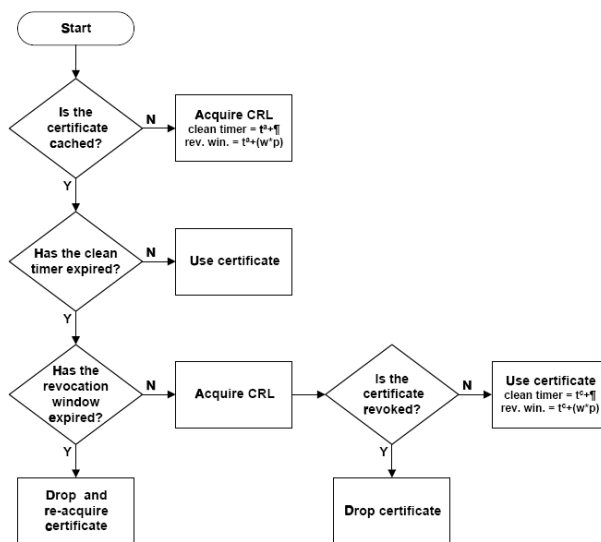


Figure 2: Verifier cache algorithm from [13]

In addition to the change at the issuer, namely the specification of the revocation window size, there is a new parameter defined by the user of a certificate. The client defines a *clean timer* for

each certificate. Put altogether, the protocol for the client is as follows (*shown in figure 2*). If a client does not have a certificate, she retrieves a fresh copy and starts her clean timer and a revocation window timer. The clean timer basically determines how fresh a certificate must be not to have to revalidate it. So if she already has the certificate and the clean timer has not expired, she simply uses the certificate without revalidating. If the clean timer has expired, she checks the revocation window timer. If the latter timer has expired, she gets a fresh certificate and resets the timers. Otherwise she retrieves the latest CRL (if she does not already have it), and checks the validity of the certificate against the CRL. If it is on the CRL, she of course drops it. If it is not on the CRL, she resets both timers and uses the certificate.

Notice that the case when the timers are always set to 0 is identical to the situation in which only fresh certificates are used. The case when the revocation window size is set to infinity is the same as using regular CRLs. Only slight modifications are needed to make this work with delta-CRLs. So this is definitely an improvement to the methods with which they compare their system. More tests would need to be performed to compare it to systems such as CRS and CRTs.

## 5 Real-time PKI Services

As Internet connectivity and accessibility have improved, people have sought real-time solutions to enhance PKI. These services can provide revocation information, offload the work of certificate validation and even be used to enforce organizational PKI policies. PKIX has proposed three such services: *Online Certificate Status Protocol* (OCSP), *Simple Certificate Validation Protocol* (SCVP) and *Data Validation and Certificate Server* (DVCS) protocols.

### 5.1 Online Certificate Status Protocol

OCSP was developed as an alternative to CRLs for the PKIX project. Its purpose was to avoid downloading long CRLs and to provide the freshest information possible about certificate revocation. An OCSP responder is a trusted server that responds to a client's request for information about the revocation status of a certificate. A positive response only means that the certificate has not been revoked. It does not imply validity, meaning the OCSP responder is not checking the signature on the certificate or its path back to a trusted root. It is not even checking that the serial number is that of an issued certificate. Obviously, the responder must be trusted. It could be trusted just to respond for that certificate if the CA issuing the certificate indicates the server as being the official responder for that certificate. A responder could also be trusted for all responses if some prior trust relationship as been established with the client. An example would be a company that has its own OCSP responder setup to do revocation checking for all employee requests.

In a way OCSP is really a step back from previously discussed certificate revocation methods. It provides shorter responses than full CRLs, but other methods such as CRS provide even shorter responses. Additionally, we are falling back on the use of a trusted third party, namely the OCSP responder. The previously discussed methods rely only on untrusted directories. While OCSP does offer the freshest information possible, CRTs can offer information nearly as fresh without use of a trusted third party. In the OCSP RFC [12], they do not have a graceful way to deal with OCSP responder key compromise. They mention that either traditional CRLs can be used for OCSP responders or their keys could be short-lived. It seems that we gain little if we are still tied to using CRLs, except in this case the list should be shorter. We feel that the better of these choices is to frequently change the OCSP keys. If archiving of OCSP requests is important, then a frequent key change could make audits more complicated. And there would still be a need to store CRLs

of OCSP keys for auditing to work. However, this is a real-time system, and the responses mean little after the fact. So auditing may not be an issue. In this case, it could be acceptable to use short-lived keys as an alternative to revoking OCSP responder keys.

An additional scaling problem comes from the fact that all OCSP responses must be signed. If they are not, someone can perform a DoS attack by faking messages that say valid certificates have been revoked. But signing every message with a public-key algorithm can overburden a server. It could lead to another type of DoS attack where a malicious user just floods the responder with requests. Caching cannot help us scale either. To prevent replay attacks, the messages must have a nonce, time-stamp or some other unique identifier. Though, if timestamps are used, a client could be configured to accept cached messages up to a certain age. But time-stamping has its own set of issues.

## 5.2 Simple Certificate Validation Protocol

SCVP is a system that allows clients to offload much of their certificate handling to a server. This can help to relieve the workload of a very low powered client, and it allows an organization to centralize PKI policies. Clients may request full validation of a certificate or just ask for construction of a certification path which it will validate itself.

SCVP servers can be trusted or untrusted. An untrusted server could supply a certification path. In [2] the authors feel that an untrusted server could also supply revocation information such as CRLs or OCSP responses. There certainly is no problem having an untrusted server give a user CRL information. We feel that it may be a little more complicated to have an untrusted server provide OCSP responses, and such a protocol must be carefully designed. Obviously, the untrusted SCVP would be giving a client information from an OCSP responder that the client trusts, though. While path construction may be trivial in single level or hierarchical PKIs, it can be quite challenging with *meshed PKIs* (collections of cross-certified CAs) or what [19] calls *bridge-connected PKIs*. Bridge-connected PKIs use *Bridge CAs* (BCAs) to connect other meshed and hierarchical PKIs. They consider SCVP servers to be a particular instance of what they call *Bridge Validation Authorities* (BVA).

A trusted SCVP server can do more. A trusted server can be used to handle almost all cryptographic work and network communication. By allowing the SCVP server to perform validation and revocation checks (if the client is interested), the client only has to send and receive one message. This could be useful for PDAs with limited wireless bandwidth and computing power (though cryptography does not take that much CPU load anymore). Even more useful may be the ability to centralize all PKI/PMI policies for an organization with an SCVP server. Using SCVP, the organization has complete control over how validation is performed. This is particularly important when SCVP servers are used as BVAs. Policies can be extremely complex and dynamic in bridge-connected PKIs, and the client software is currently not intelligent enough to interpret and process all of those policies.

The authors do note a few important issues. First, a trusted SCVP server is trusted as much as a root CA. So the keys must be strong and protected carefully. Clearly, compromise of the key is detrimental and could result in a client accepting ANY bogus certificate. Also, it is recommend that the client use an unpredictable sequence of identifiers for requests so that it does not fall prey to replay attacks. Lastly, they point out that policy information requests and responses are not signed, and hence vulnerable to man-in-the-middle attacks.

Our biggest problem is that the servers are very heavily loaded, making all of the cryptographic workload even more unbalanced. This makes the system even more unscalable. With desktops or laptops, the client usually has more free CPU time than the server, and moving the burden to

the server exacerbates the situation. So if the client is not a small wireless device with limited bandwidth, the only use we see is in the centralized PKI/PMI policy making. This can be quite an advantage in many situations, though. This actually helps scaling with bridge-connected PKIs because it provides quick updates of complex sets of policies and may be necessary since most clients are not intelligent enough to interpret and act on those policies. Here an SCVP server acting as a BVA might not be heavily burdened if it is just set to deal with certificates for other domains that the client does not understand.

### 5.3 Data Validation Certificate Server Protocols

DVCS is not a replacement for CRLs or OCSP. The purpose is to extend functionality. In fact, DVCS could not replace CRLs in a large open environment due to scalability issues. A DVCS is like a notary public. It is used to bind a time to a particular event, such as the signing of a document. A DVCS issues a Data Validation Certificate (DVC) signing that something happened or was valid at a given time. More specifically, it provides the following services.

*Certification of Possession of Data* is a DVC that states a requester possessed data at time  $x$ . This is essentially a time stamp by a trusted third party, namely the DVCS. *Certification of Claim of Possession of Data* is almost the same, except that the requester only shows the DVCS a hash of the data. This is useful if the data needs to be kept private. Again, this is basically just a time stamping service. *Validation of Digitally Signed Documents* is a service that checks signatures on a document, verifies that they are good at a particular time, and signs a DVC stating this fact. *Validation of Public Key Certificates* is the same except that the DVC is validating that a PKC is good at a particular time. This implies that the DVCS checked the path to a root CA, as well as the certificate in question.

The main benefit of these services is non-repudiation and extension of signature validity. By having a DVC, an auditor can see that the document signature was valid at the time DVC was issued. It doesn't matter whether the signature key has now been expired or revoked. Without this service a signed document must still be reliably time stamped, and an auditor would have to check archives of CRLs to determine the validity of the key (and others in the verification path) at the time of the original signature. But now the signature is valid until the DVCS's key expires. However, this can be extended by the DVCS issuing a new DVC before its key expires.

As the authors of [3] point out, use of a DVCS would be helpful when performing a transaction involving large sums of money. Not only does it check validity of the key for a client (using OCSP, CRLs or other methods), it provides a DVC which can be used for non-repudiation if needed. However, there is a lot of computation and communication that the DVCS provides for the client by doing these checks. So we would see DVCS use being a pay service, and likely not to be needed all the time. This is good since it would be hard to scale given the server burden. Another use might be for a corporation to setup a DVCS server that employees are required to use. This would create an audit trail, and it would allow the company to set strict policies on verification of certificates via the DVCS server.

It should be noted that the client still does have the responsibility of checking the validity of DVCS server certificates through traditional methods. In a corporate situation, the client could rely on the fact of being notified immediately of a compromised DVCS key. Other methods may be to use OCSP or CRLs for DVCS key revocation. Using either method, such a compromise is very damaging since it invalidates all the previously issued DVCs with that key. If a DVC is being used to extend the lifetime of a signature and the DVC is compromised, the signature is now useless. Redundancy, such as the use of two DVCSs at all times could help, but it is not a solution that helps the scaling issues. Strong keys and serious methods to protect them are certainly in order.

## 6 Military Considerations

### 6.1 Enrollment and Distribution

Military instances of PKI have several unique factors to consider that will effect the scalability and usefulness of the system. At the highest level, one must decide on the basic structure. It is favorable that PKIs tend to be hierarchal, much like the military structure. The armed forces can use the existing structure and overlay the PKI upon it. For example, there could be a root CA that controls policy and distributes a few certificates to all five branches of the military. Then the CAs for each branch could determine another level of CAs underneath them. It would be unrealistic to assume, for example, that there is only one CA to issue certificates for the entire Navy. Another consideration is whether Registration Authorities (RAs) should be used. While the CAs may well be capable of handling the computational load of certificate generation and revocation with only three levels in the hierarchy, physical distance might become a larger problem with distribution and enrollment. One certainly would not expect that an officer or a soldier would go back to their home base to prove their identity in person every time they needed a new key. So mobile RAs could be setup anywhere troops are deployed to verify identities, generate keys, send signed certificate requests to CAs, and distribute certificates.

Another major architectural decision is whether attribute certificates are used or roles are part of the identity in a PKC. Being that soldiers frequently change roles -in fact, each mission may be considered a new role- it would be useful to have different certificates per role. If the role is part of the identity, identity certificates, which require a user to go to an RA or CA, would need to be issued very frequently. It is easier to distribute an attribute certificate since the user doesn't need to receive any private data, such as a private key, with it. Also, the attribute certificates can be very short term to prevent its use past the short lifetime of a role or mission. Being that they have such a short lifespan, it should be uncommon for them to be revoked.

At a more physical level, the military has special requirements. If smart cards are used, they must be physically protected from the elements and potential abuse. These are sensitive pieces of technology that would need to be physically protected by some sort of case, at least from ElectroStatic Discharge (ESD) if nothing else. An alternative may be to use a rugged smart card, but this solution would be expensive not only because of increased manufacturing costs but because there is little competition in that market. A generic smart card costs on the order of tens of cents and a protective case may be on the order of a couple dollars or less. Special rugged smart cards could easily cost much more. Additionally, smart cards require smart card readers. All equipment for communication would have to be replaced by newer equipment that utilizes smart cards. A software solution may be a more feasible choice. Such a solution could take the form of traditional credential wallets, such as the Verisign Roaming Service or the NSD Security Roaming Solution. In this situation the soldier could download her credentials into the communications hardware for use, after which it would be wiped clean. Not only does this require less special hardware to carry and buy, but it solves the problems of distribution and enrollment. Using this online method, a soldier can always ask for and receive new credentials without physically going anywhere. Additionally, the private credentials are not permanently stored on the soldier. The biggest problem is of course that it does require online use. In this system, orders could not be signed or decrypted offline.

### 6.2 Interoperation

Secure communication solutions can be difficult to deploy and manage in a multi-organizational effort. Different organizations use different PKIs, and it is often impractical to create a unique PKI

for the mission because every single user in each organization must then apply for a certificate that is valid just for the one mission.

As explained in section 5.2, we can form a bridge-connected PKI having several PKIs connected by a Bridge CA (BCA). SCVP, or more generally Bridge Validation Authorities (BVAs), can be used for smoother integration of different administrative domains within a bridge-connected PKI. BVAs facilitate certification path discovery and can respond to clients about the status of certificates in domains other than their own. Policies for cross-certification and validation are enforced by BCAs and BVAs; so any new policy specifications or changes need to be done at that level and do not require cooperation of the involved PKIs. This is an important property that is needed to add or remove organizations. Typically what happens is an organization in this larger bridge-connected PKI will have its own BVA to control how the users of that domain deal with certificates in other domains. So the BVAs control policy, and the BCAs simply provide connectivity so that certification paths can be constructed.

Bridge-connected PKIs solve the problem of interconnecting PKIs having similar structure, for example CA-based PKIs. If different implementations of PKIs need to communicate, we need a more sophisticated solution. We define a *Translator* as an intermediary entity helping those different PKIs communicate with each other. If a user of a PKIX-like PKI wants to talk with another user in a PGP-like PKI, the translator will contain a set of rules/policies to transform the information from the PKIX certificate to something that can be understood by the PGP user. For example, the translator can get the public key out of the certificate, sign it with its own private key, and send the result in a PGP message. The PGP user will trust the key since it comes signed by her translator.

### 6.3 Offline Operations

Common certificate revocation methods require online operations. OCSP is a real-time service, and as such it obviously requires users to be online. CRLs could work offline if the user were able to carry the latest CRLs with her. That way she could prove to any other user that her certificate has not been revoked. However, CRLs can easily be megabytes in size and hence not practical to store on mediums such as smart cards (Although PDAs may be suitable for the task).

As we noted in chapter 4, CRS and CRTs provide more succinct proof of validity that can be carried on smart cards. Moreover, CRS provides the extra benefit of short term validation of certificates. With a period of one day, the CRS signature is only good for that one day. So a captured CRS signature along with a compromised key would only be useful for a short period of time. Of course, this length of time must be balanced with how often a user will need to go online to get the new validations. While it may be quite secure to have the CRS statements valid for only an hour, it is impractical to require users to go online to get the latest signatures every hour.

While it is more natural to talk about this period for CRS, since it is an integral part of the CRS protocol and algorithms, short periods of validity can be used for any revocation system. CRTs could be given short lifetimes either by putting an end-of-life time with the CRT messages, or by indicating when the next update will be issued. Such information must be signed along with the root of the CRT, of course.

### 6.4 Fault Tolerance

In most military applications, fault tolerance is necessary. Certificate revocations systems are no exception. While replication of directory services can increase the fault tolerance of such systems, some situations require more robust solutions. In [17], Wright et al. develop a fault tolerant network

to distribute information while balancing the load on each node. While they focus on using it for certificate revocation, it could be used for any type of information. They accomplish this through the use of what are called depender graphs, a special type of Directed Acyclic Graph (DAG). They have a parameter  $k$  which specifies the minimum number of nodes that must be disabled before failure can occur. This parameter is inversely related to the load put on each node in the graph.

The original design has a single root which pushes information through the graph to all interested parties. Wright et al. [17] note that by using threshold certificates and multiple roots, even the root can be made tolerant to failures. Some real advantages compared to many distributed systems is that the load can be kept balanced by the parameter  $k$ , and this system does not require knowledge of global state. Each node only maintains information about  $k$  parents and up to  $k$  children. Also, the system is general enough to support more than just CRLs. It could support PGP revocations and CRS messages as well. However, this system would not be practical for PGP if users wanted to send revocation information directly into the system since they would need to become roots of their own depender graphs. If all PGP users did this, each would need to join depender graphs for every key they are interested in hearing revocation information about. Thus it would be more practical for PGP users to send revocation information to a common directory which would push this revocation information to users. However, in military scenarios PGP may simply not be viable because only the key owner can revoke a key, and only if she is alive.

One limiting factor to scalability is the fact that nodes must keep all messages they receive to pass to new members. This could be a lot of state and overwhelm new users or nodes that left the network for a while. However, in the case of certificate revocations, they need only provide information about non-expired certificates. If CRLs are being used, only the most recent need to be stored and any delta-CRLs since the most recent full CRL. An assumption is that all nodes will pass on information whether they are interested or not, and they will do so in a timely fashion. The fault tolerance can make up for some lazy nodes, but there is still some sort of trust a user must have in its parent nodes.

Wright et al. briefly discuss some issues related to reconfiguration after failures, but they do not present a solution of their own. Before a system like this is used in production, we feel that this issue must be resolved.

## 7 Conclusions

We have briefly shown the need for PKIs and discussed some of the popular implementations, namely PKIX, SPKI and PGP. In all of these systems there is a need to perform both efficient enrollment and revocation. In chapter 3, we looked at how FreeICP and Self-Assembling PKIs try to address the scaling issues related with enrollment. We also looked at how the Canadian government has tried to use PKI while preserving privacy as best as possible. In chapter 4 we examined some of the more common certificate revocation methods. We looked at traditional CRLs and improvements upon them but noted that they still do not scale well. While delta CRLs do scale better, they do not provide succinct proof of validity that a user can carry with her. Both CRS and CRTs provide much more scalable solutions to certificate revocation and also provide succinct proofs of validity. These two solutions do differ in how they balance the amount of communication between the directory and CA with the amount of communication between the directory and the end users, but they are both very balanced compared to traditional CRLs.

In chapter 5 we looked closer at some of the new real-time PKI services such as OCSP, SCVP and DVCS. These services offer everything from real-time certificate status checking to complete certificate validation and verification. SCVP even allows organizations to create central points

of management for all certificate handling and PKI policy enforcement. Lastly, in chapter 6 we discussed how some of the technologies we have looked at can be useful to military applications. We saw how SCVP and related technologies can help integrate the PKIs of a newly formed coalition. We also noted how CRS and CRTs can help users who are not able to be online as much. And lastly, we discussed dependency graphs and how they create redundant channels of communication that can be utilized for certificate revocation or update information.

There are several topics that we intend to examine in the future. Some are other certificate revocation systems, including windowed certificate systems. We will be looking more closely at role-based cryptosystems as a possible solution for one-to-many communications. There are many more technologies aimed at secure information sharing among coalitions. One such technology we would also like to examine is translators that integrate the use of different kinds of PKIs. Traditionally, PGP users cannot communicate with PKIX users. However, translators sit between the different systems and can enable communication between SPKI, PKIX and PGP users almost seamlessly. Additionally, we are interested in recent work that models PKIs and especially certificate revocation systems. A formal model in which to place different systems will help us compare different PKIs in a consistent manner.

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